

Typed Meta-Interpretive Learning for Proof Strategies

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Motivation

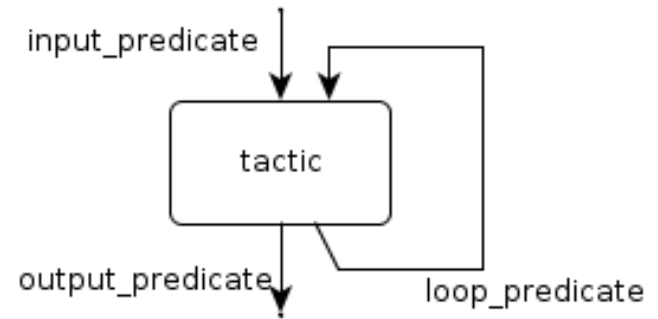
- Use of theorem proving software increasingly widespread in industry
- Manual intervention often required; time-consuming and requires a skilled user
- “Stuck” proofs often cluster into groups which can be addressed using the same strategy

Motivation

- We demonstrate that *meta-interpretive learning* (MIL) can be used to learn these strategies in order to re-apply them elsewhere
- These strategies have a high degree of branching and thus a large search space
- We introduce *typed MIL* and hypothesise:
“*Typed MIL learns more deterministic proof strategies than untyped MIL*”

PSGraph

- Graphical representation of proof strategies [2]
- Nodes represent tactics, predicates on wires direct goals between nodes and ensure correctness



An example PSGraph with edges labelled by predicates

Metagol

- We use the *Metagol* [3] implementation of MIL
- Encode tactics and goal data as background information
- Successful proof evaluations given as positive examples
- Metarules applied to background to find a definition which satisfies examples

Metagol

- Metagol allows *predicate invention*, i.e. missing definitions can be found using available background information
- We can use this to find wire predicate definitions:

```
strategy(A,B) :- strategy_1(A), tactic(A,B).
```

```
    strategy_1(A) :- has_symbol(A,C).
```

Metagol

- Two separate learning problems: structure and conditions
- To learn both we must either restrict wire predicates or learn a strategy with a high degree of branching
- To resolve this we introduce *typed MIL*

Typed MIL

- All predicates and arguments are “tagged” with a constant representing their type:

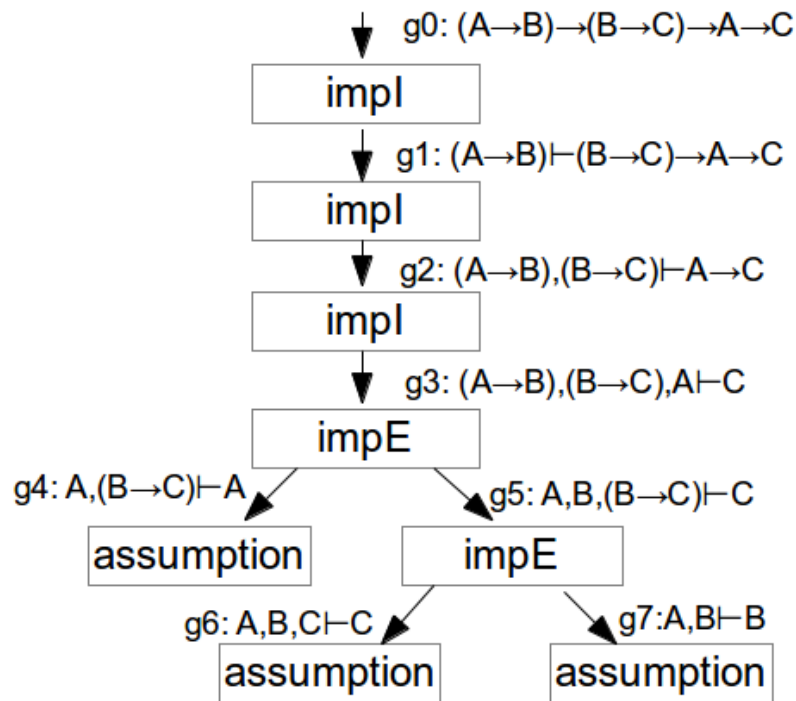
$$P(X, Y) \longrightarrow P:t_1(X:t_2, Y:t_3) .$$

- To work in Metagol we treat the predicate type as an extra argument:

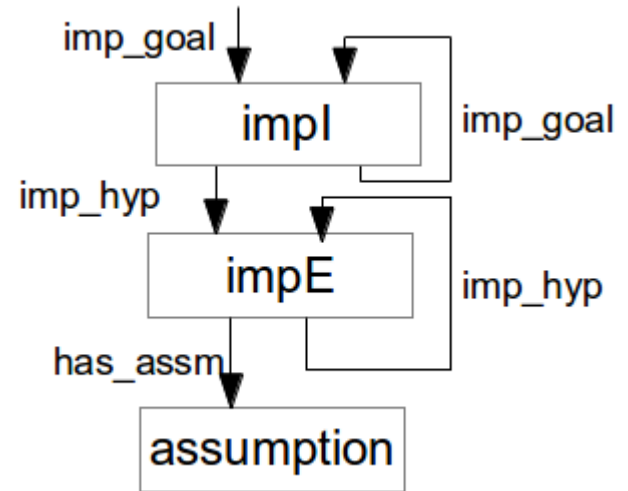
$$P:t_1(X, Y) \longrightarrow P(t_1, X, Y) .$$

An Example Strategy

$$(A \rightarrow B) \rightarrow (B \rightarrow C) \rightarrow A \rightarrow C$$



Proof Tree



PSGraph

Learned Definitions

Untyped Strategy

```
strat_i(A,B) :- erule_impE(A,C), assumption(C,B).  
strat_i(A,B) :- erule_impE(A,C), strat_i(C,B).  
strat_i(A,B) :- rule_impI(A,C), strat_i(C,B).
```

Typed Strategy

```
strat_i(psgraph,A,B) :-  
    assm_type(wpred,A), assumption(tactic,A,B).  
strat_i(psgraph,A,B) :-  
    strat_i_1(psgraph,A,C), strat_i(psgraph,C,B).  
strat_i_1(psgraph,A,B) :-  
    impE_type(wpred,A), erule_impE(tactic,A,B).  
strat_i_1(psgraph,A,B) :-  
    impI_type(wpred,A), rule_impI(tactic,A,B).
```

Results

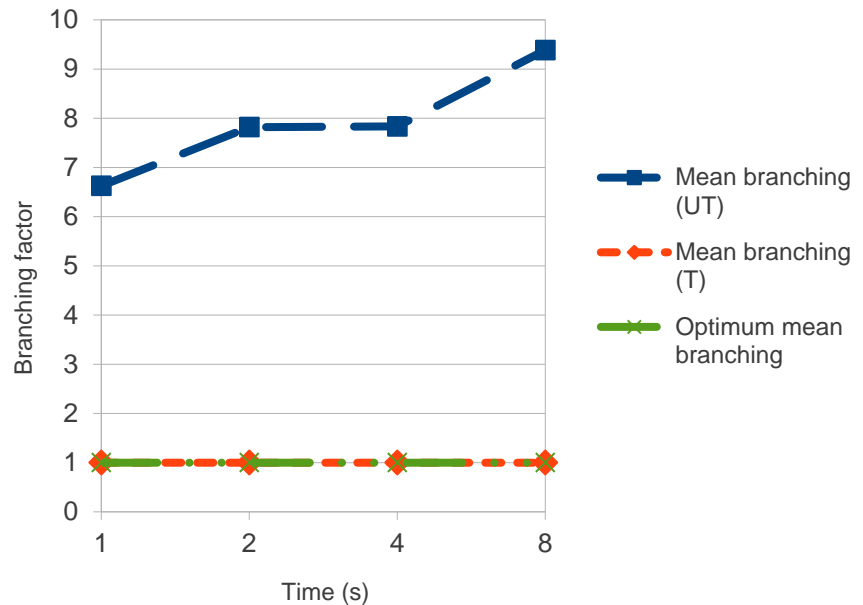
- Learned strategies from 15 propositional logic proofs
- Metagol run for 1, 2, 4 and 8 seconds

	Successes	Mean nodes	Mean clauses	Mean branches	Mean evaluations
Untyped	13	4	3	9	1
Typed	7	4	4	1	2

Average learning results across all experiments

Results

- Consider branching factor σ for each strategy.
- Large σ indicates large search space
- Untyped: $\sigma > 1$ and increases
- Typed: $\sigma = 1$ and constant



Mean branching factor of strategies learned using untyped (UT) and typed (T) MIL compared to the optimum branching factor

Conclusions

- Typed strategies have less branching than untyped strategies
- Thus typed strategies are more deterministic
- Initial results using this approach have been very promising

Further Work

- Developing strategies from multiple proofs – already underway. Will use examples from group theory.
- Moving on to larger, more complex proofs, e.g. rippling [1]. A simplified version can be learned with untyped MIL, can we learn a full version with types?

References

- [1] A. Bundy, *Rippling: meta-level guidance for mathematical reasoning*, Vol. 56. Cambridge University Press, 2005.
- [2] G. Grov, A. Kissinger and Y. Lin, *A graphical language for proof strategies*, in *LPAR*, volum 8312 of *LNCS*, pages 324-339. Springer, 2013
- [3] D. Lin, E. Dechter, K. Ellis, J. B. Tenenbaum and S. H. Muggleton, *Bias reformation for one-shot function induction*, in *ECAI*, pages 525-530, 2014

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