

# Languages and states (another view of “Why”)

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# Menu

- one technical postscript
  - ... to other talks on AI4FM project
  - recall: my *personal* focus on CxC
- a bunch of (polemic) points
  - if time permits
  - “an idealist’s response to justifiable criticism”

# “Models of Why”

- (as Leo said)
  - we concocted this to indicate what might be extractable from proof (attempt)s
  - of course, we have yet to perfect it!
  - nor have we published it beyond project
  - have a “state” ( $\Sigma$ )

# Designing a language: advice

- postpone concrete syntax
  - as long as possible
- even, postpone abstract syntax
- focus on the underlying state  $\Sigma$ 
  - (before thinking about (SOS) rules)

# Isabelle (e.g.)

- $\Sigma$ 
  - pretty flat!
  - current goal list
  - ...

# Change of view

- (recent)
- we are designing the state of **the language** that helps express learning and replay

## ProofProcess model

- Intent
- Proof features

- Attempts
- Granularity
- Structure

Abstract  
term data

- Proof scripts
- Proof history

## Create ProofProcess Model

Ask expert what's  
happening

Infer proof  
process

“Clippy”

Proof replay

App X.. ?

Apps

## ProofProcess Integration

PP: Isabelle

PP: Z/Eves

PP: Rodin

PP: ...?

Theorem Proving  
System



# Architecture

# Models of why

as  $\Sigma$

- (collection of *Theories*)
  - “inheritance” links
- defines *Operators*
- and *Conjectures*



# Conjectures

- can have *Proof*
  - yes, the whole thing!
  - structured
- but can also be *Axiom, Trusted, Tool*
- (or be incomplete)
- can also have *Disproof*
  - Aaron point about negative information

# *Conjectures* might also have

- *Shape*
- *Uses*
  - set of *Clue*
- etc., etc.

# Details are unimportant today (but “vital” later)

- crucial: we fix the state of this language *before* worrying about its statements
- from this viewpoint:
  - we are trying to “parse”/prompt a new attempt
  - match to previous “programs” (graph matching)
  - clear role for machine learning

# Further language issues

- the language has to harness parallelism
- “non-procedural”?
- the form of this language might be unconventional
- some of activity at model (pre PO) stage

# Polemics

- there will be a horizon for *any* TP ideas
- extra model layers just to reduce TP task?
- we are trying to harness an extra resource
  - the results of an expert doing one proof
- there is evidence for “families”
- lessons from “mural” (cf. Ursula’s archeology)

The  $\mu$ ral Store

<table border="1" style="width: 100%; border-collapse: collapse;"> <tr><td>developments</td></tr> <tr><td>lift control</td></tr> <tr><td>lookup</td></tr> <tr><td>reactor</td></tr> </table>	developments	lift control	lookup	reactor		<table border="1" style="width: 100%; border-collapse: collapse;"> <tr><th>theories</th></tr> <tr><td>Boolean</td></tr> <tr><td>Cartesian Product</td></tr> <tr><td>Finite Maps</td></tr> <tr><td>Finite Sequences</td></tr> <tr><td>Finite Sequences</td></tr> <tr><td>Finite Set Theory</td></tr> <tr><td>Integers</td></tr> <tr><td>lookup lev 1 reified</td></tr> <tr><td>lookup lev 1 theory</td></tr> <tr><td>lookup lev 2 theory</td></tr> <tr><td>LPF +</td></tr> <tr><td>Natural Numbers</td></tr> <tr><td>Non-Disjoint Union</td></tr> <tr><td>Predicate LPF</td></tr> <tr><td>Propositional LPF</td></tr> <tr><td>reactor reified by</td></tr> <tr><td>reactor theory</td></tr> <tr><td>reactor1 theory</td></tr> <tr><td>Typing Assertion</td></tr> <tr><td>VDM BASE</td></tr> <tr><td>VDM LOGIC &amp; D</td></tr> <tr><td>Weak Equality</td></tr> </table>	theories	Boolean	Cartesian Product	Finite Maps	Finite Sequences	Finite Sequences	Finite Set Theory	Integers	lookup lev 1 reified	lookup lev 1 theory	lookup lev 2 theory	LPF +	Natural Numbers	Non-Disjoint Union	Predicate LPF	Propositional LPF	reactor reified by	reactor theory	reactor1 theory	Typing Assertion	VDM BASE	VDM LOGIC & D	Weak Equality
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Proof for  $\forall$ -E

rule:	$\forall$ -E	attempts
theory:	Predicate LPF	main proof
marked items:	<input type="checkbox"/> consistent <input type="checkbox"/> complete <input type="checkbox"/> not assumed	<input type="button" value="clear"/>
<input type="button" value="tactic tool"/> <input type="button" value="justif tool"/>		

  

main proof

```

h1  $\forall x : X . ( P [ x ] )$ 
h2  $( a : X )$ 
1  $( \neg \exists x : X . ( \neg ( P [ x ] ) ) )$ 
2  $( \neg ( \neg ( P [ a ] ) ) )$ 
c  $( P [ a ] )$ 

```

Proof

spawn proof

- show unproven rules used
- clean up proof
- modify
- make proof assumed
- redraw

unfolding from h1  
by  $\neg\exists$ -E on [1, h2]; []

by  $\neg\neg$ -E on [2]; []

  

Predicate LPF

parents
Typing Assertion
rule groupings
$\forall$ -E : Rule
$\exists\forall\rightarrow\forall\exists$ : Rule
tactics
StripExistential
StripQuantifiers
StripUniversal
oracles

  

$\forall$  : Binder

```

{
  (  $\neg \exists x : A . ( \neg ( P [ x ] ) )$  )
}

```

accept

  

$\exists\forall\rightarrow\forall\exists$  : Rule

```

{
  (  $\exists x : X . \forall y : Y . ( E [ x , y ] )$  )
}

```

accept

  

$\forall\wedge$ -dist-expand  
 $\forall\rightarrow\exists$  : Rule  
 $\forall\rightarrow\neg\exists$ -deM : Rule  
 $\forall\forall$ -comm : Rule  
 $\forall\wedge$ -dist-contrac

# “mural”

- interesting GUI experiment
  - but no decision procedures!
- built from formal spec
  - which was maintained!!
- [JLM91] available as:
  - [homepages.cs.ncl.ac.uk/cliff.jones/ftp-stuff/mural.pdf](http://homepages.cs.ncl.ac.uk/cliff.jones/ftp-stuff/mural.pdf)