



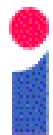
Towards a Strategy Language for Describing Proofs

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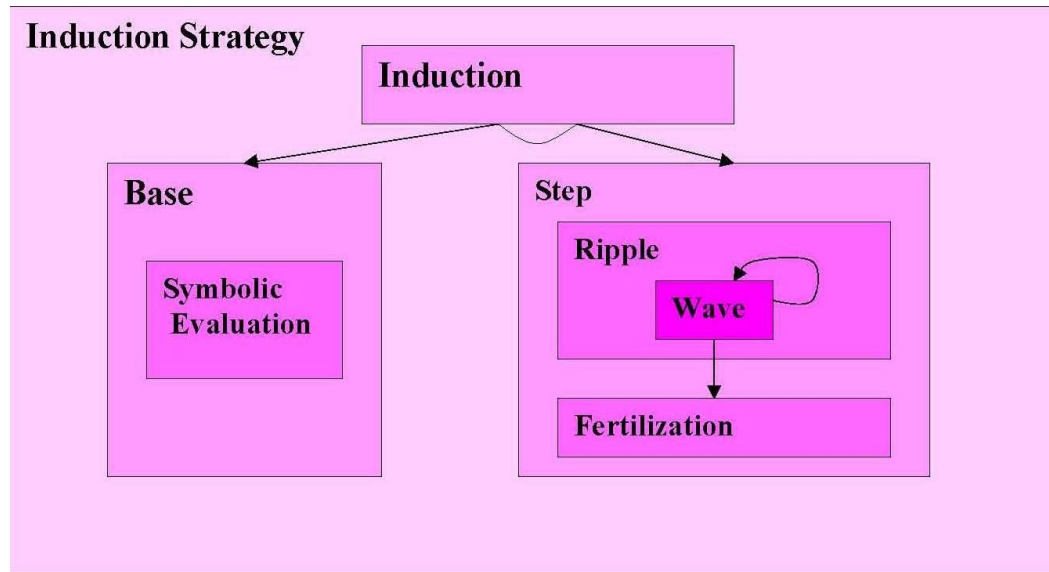
AI4FM Kick-off Meeting



Outline

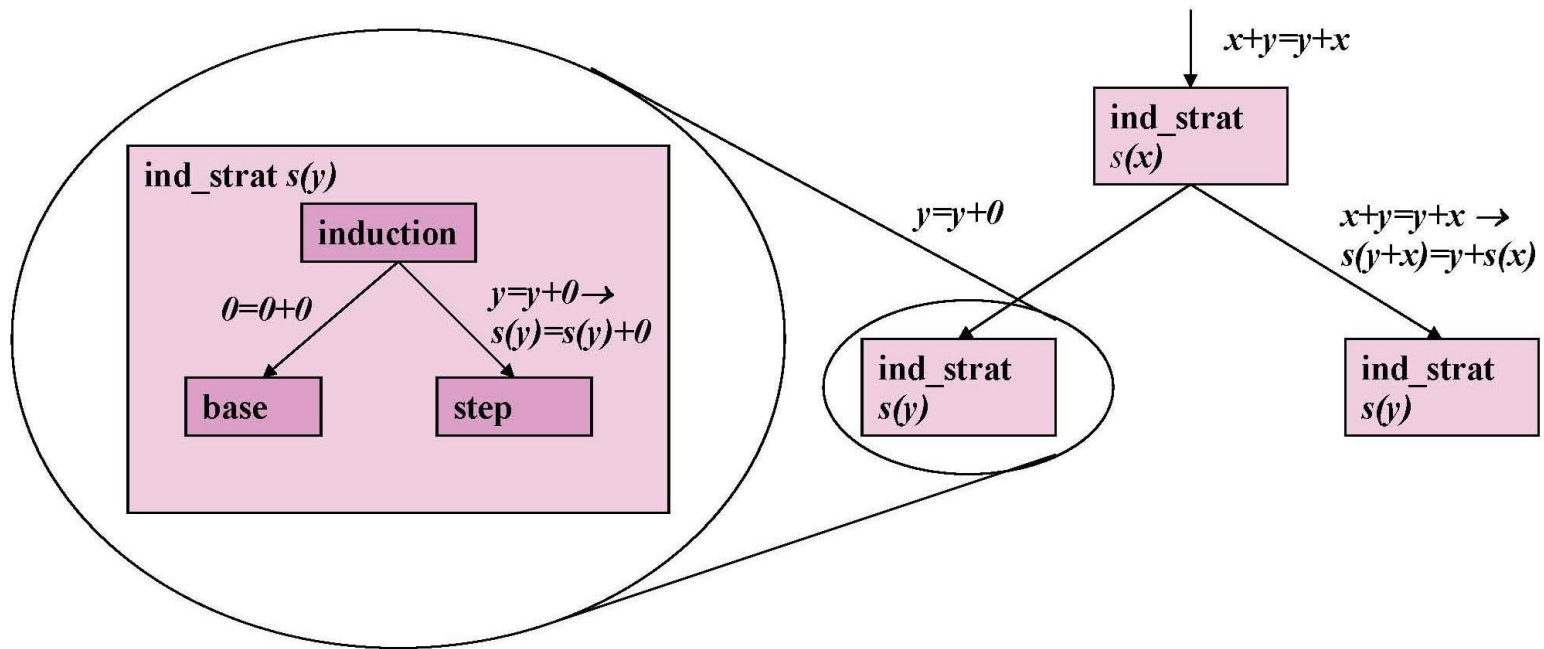
- Introduction to hiproofs
- Hiproofs as candidate strategy language
- Can describe search spaces too
- Consider alternative 'ACL2' methodology

What are Hiproofs?



- **Higraphs**: graphs whose nodes are higraphs.
- **Hiproofs**: use of higraphs to represent proofs.
 - <http://dream.inf.ed.ac.uk/projects/higraphs/>
- **Nodes** represent tactics; **arcs** pass subgoals.

Represent Proof at Different Levels of Detail



Zoom in and out of proof

Motivation

- Can describe transferable, high-level structure of proofs.
- Can parse source proof to level of detail permitted by known tactics.
 - 'black box' unparsable sub-proofs.
- Can generate target proof to varying levels of detail.
 - If low-level fails, try higher-level.
- Can describe partial proofs and search spaces too.

History

- Bundy & Ireland used box diagrams informally to describe proof plans.
- Denney, Power and Tourlas developed hiproof theory.
- Various hiproof viewers developed for Dixon's IsaPlanner.
- Aspinall, Denney & Lüth developed Hitac language and semantics.
- Plan to use hiproof proof constructor.

Bluffer's Guide to Rippling

- Use when 'given' embeds in 'goal'
 - e.g., inductive proofs
- Annotate to distinguish embedding from rest
 - language of wave-fronts and skeletons
- Manipulate goal so that embedding is complete sub-term
 - use annotated rewrite rules
- Use given to rewrite this sub-term,
 - e.g., to true.

Example of Rippling

Wave-rules

$$\begin{array}{l}
 H\# \boxed{T} @ L \Rightarrow H\# \boxed{T @ L} \quad (1) \\
 \boxed{X_1\# X_2} = \boxed{Y_1\# Y_2} \quad \Rightarrow \quad \boxed{X_1 = Y_1 \wedge X_2 = Y_2} \quad (2)
 \end{array}$$

Rippling

$$t @ (Y @ Z) = (t @ Y) @ Z$$

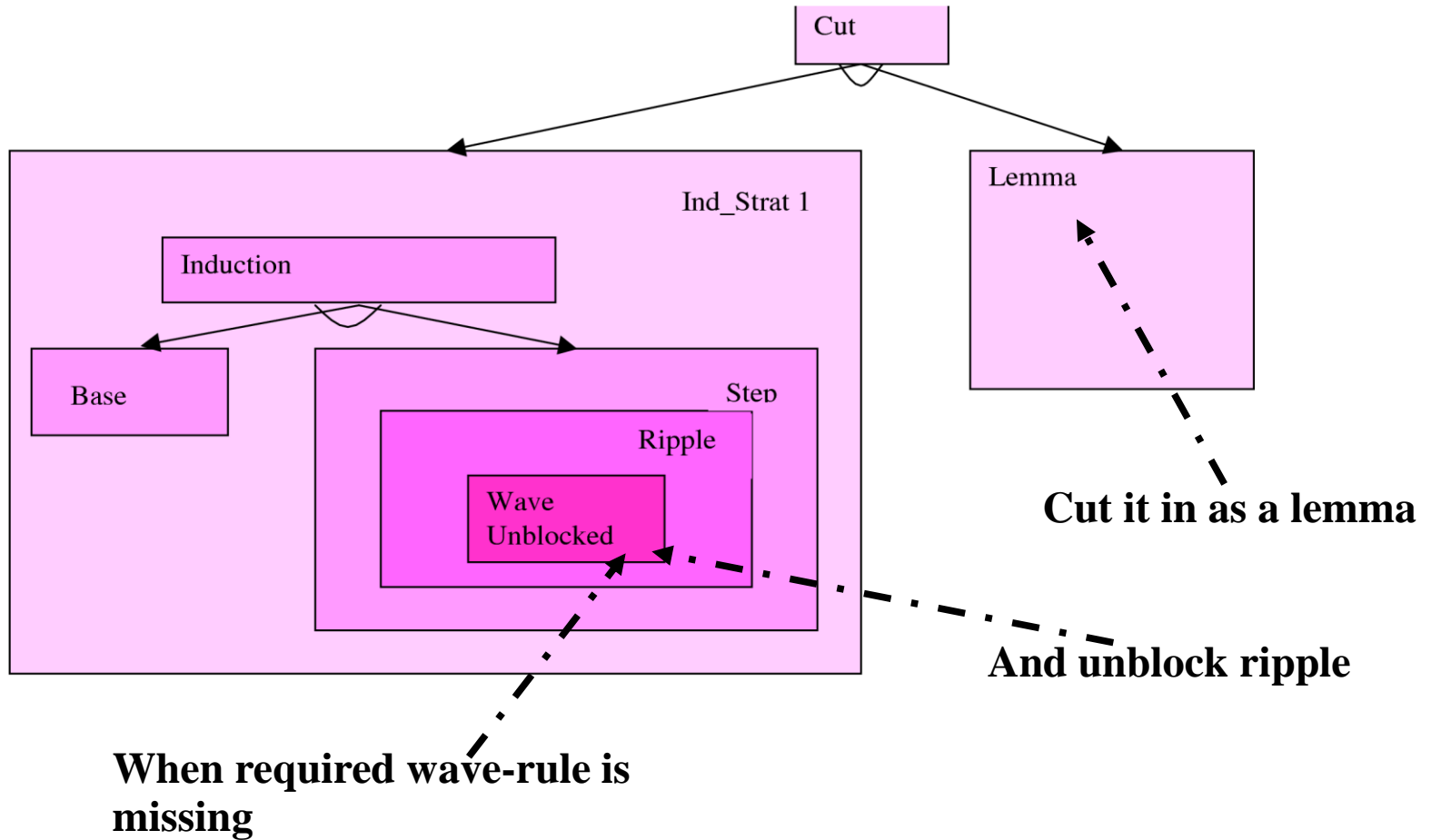
$$\Rightarrow h\# \boxed{t} (y @ z) = (h\# \boxed{t} @ z)$$

$$\text{by (1)\&(2)} \Rightarrow h\# \boxed{t @ (y @ z)} = h\# \boxed{t @ y} @ z$$

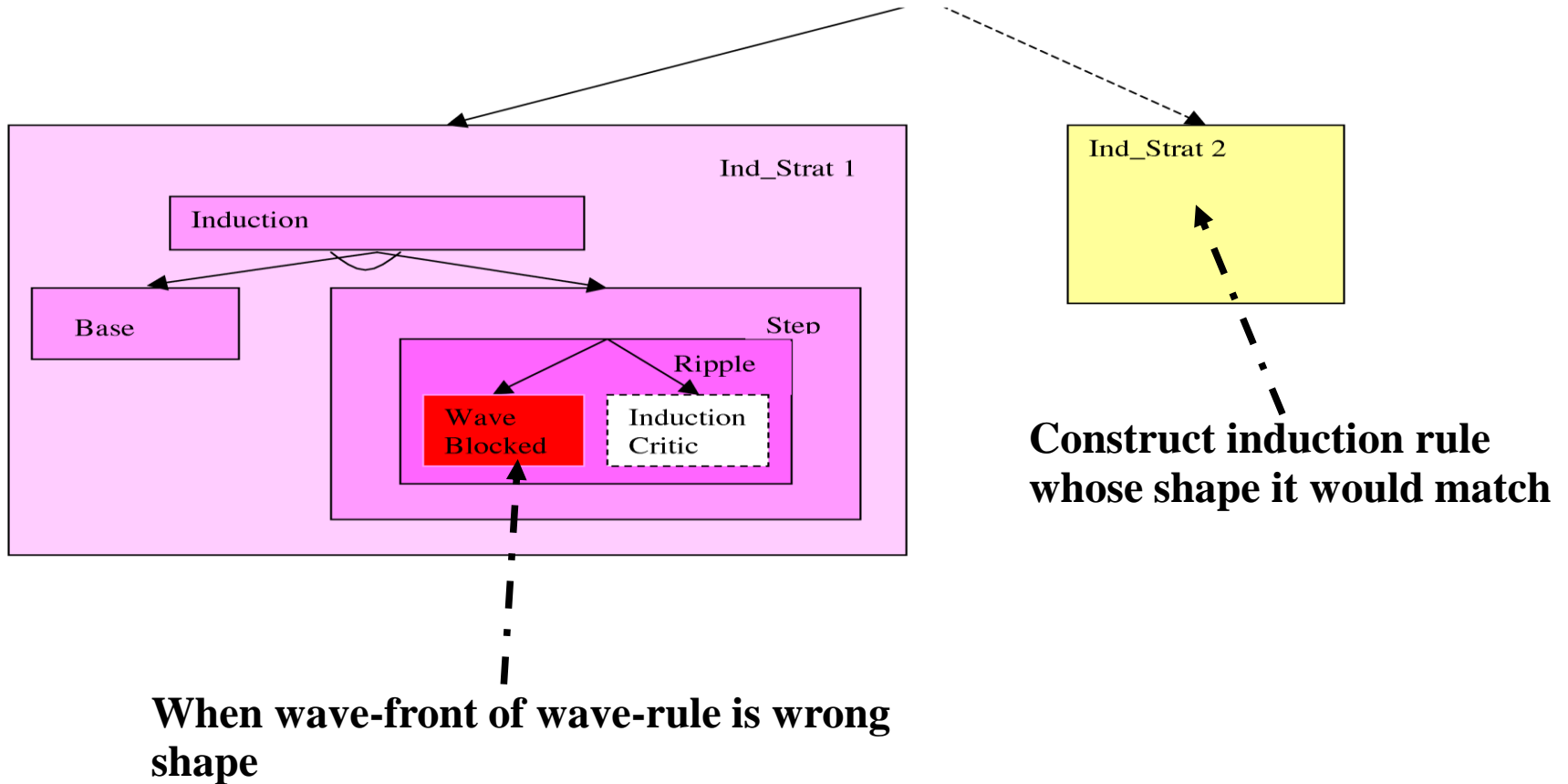
$$\text{by (1)} \Rightarrow h\# \boxed{t @ (y @ z)} = h\# (t @ y) @ z$$

$$\text{by (2)} \Rightarrow h = h \wedge \boxed{t @ (y @ z) = (t @ y) @ z}$$

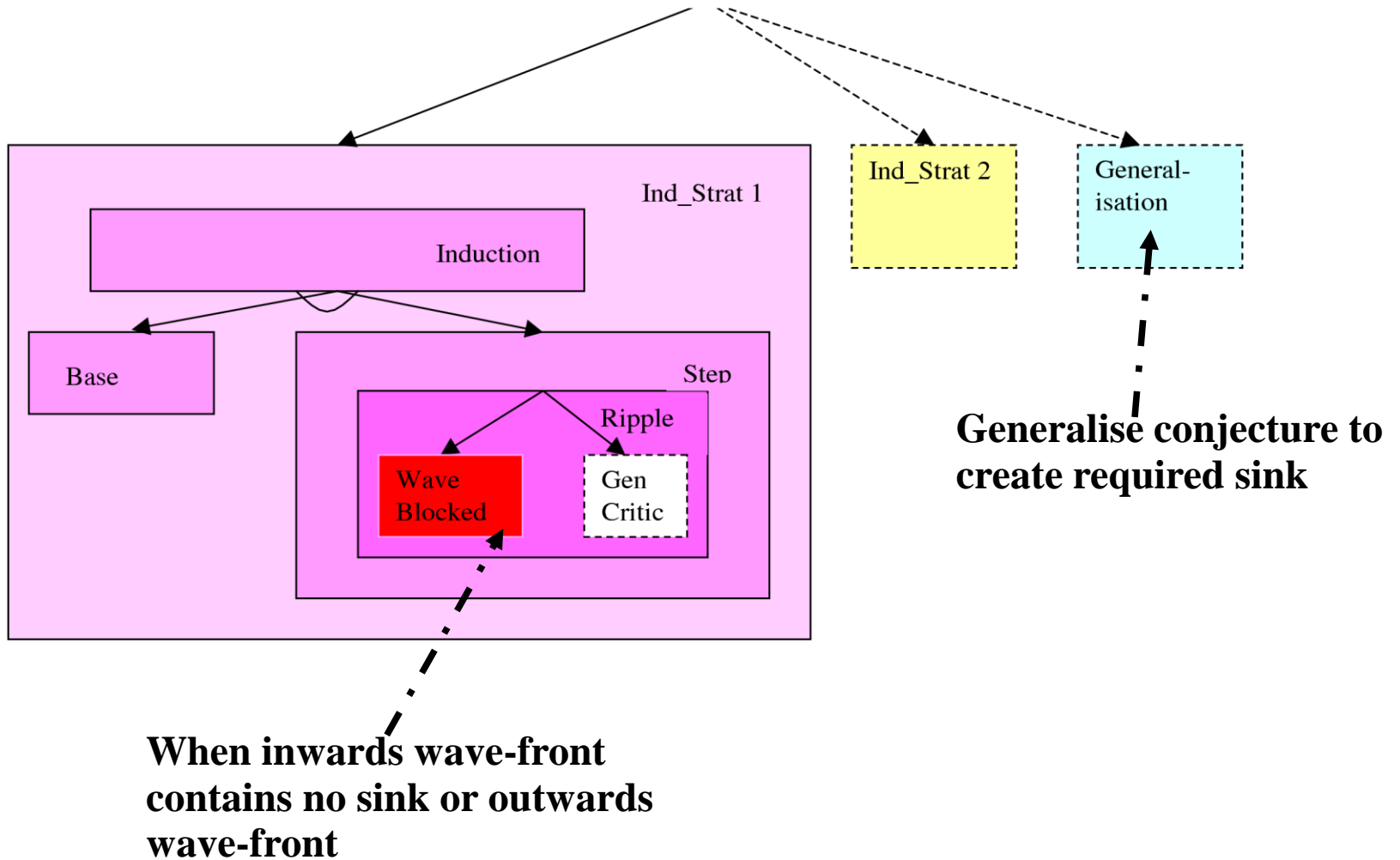
Introduce Intermediate Lemma



Change Induction Rule



Generalise Conjecture



Alternative *Unstructured* Option

- Follow 'ACL2' methodology
- 'Strategy' as provision of key lemmas
- Identify lemmas provided in source proof
 - Maybe of 'toy' conjecture whose proof fails in the 'same' way.
- Construct corresponding lemma
 - Maybe the same lemma or similar one
- Prove and add lemma – recurse
 - Maybe to particular rule set

Conclusion

- Hiproof as 'strategy language' for describing proof structure
- Parse source proof then use to guide target proof
- Limited by available tactics
 - but 'black box' unparsable sub-proofs
- Alternative is 'ACL2' methodology
 - no structure – identify and add lemmas