



Verifying Z Specifications with Z/Eves

AI4FM kick-off meeting talk

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Overview

- Mechanising Z specifications
- Z/Eves theorem prover
- Verification activities
- Outcomes & conclusions

Mechanising Z specifications

- Grand Challenge (GC) in Verified Software
 - Pilot projects: Mondex, POSIX file-store, verified OS kernels
- Existing Z specifications
- Manual or no formal proofs
- Mechanising the specification
 - Perform/verify proofs
 - Component reusability
 - Model extensibility

Z/Eves theorem prover (I)

- Successfully used for Z specification and mechanisation in GC pilot projects
- Z syntax & type checking
- User friendly/easy learning curve
 - Strict & limited number of proof tactics
- EVES proof system with Z extensions
- Step-by-step and push-button proof approaches

Z/Eves theorem prover (II)

- Supports sizeable specifications
 - Mondex benchmarks:
 - ~200 paragraphs and >300 proofs
 - >4500 proof commands
 - ~1 hour to complete proofs (laptop, 2006)
- *No active development & distribution*

Z notation

- Based on Zermelo-Fraenkel set theory and first-order predicate logic
- Z schema
 - Named record with invariants
 - Can be included in other schemas
- State & operations specification style
 - Operation schema – before and after state with transition invariants
 - Supported by Z/Eves

Verification with Z/Eves (I)

- Syntax & type checking automatic
- Domain checking
 - Z notation does not impose undefinedness treatment
 - Well-defined partial functions (Z/Eves)
 - Generates domain check proof obligations (PO) automatically
 - Conditional expressions (J. McCarthy):

$$a \wedge b \hat{=} \mathbf{if} \ a \ \mathbf{then} \ b \ \mathbf{else} \ \mathit{false}$$

Verification with Z/Eves (II)

- Inconsistency checking
 - Checks the feasibility of the specified model
 - Axiomatic definitions

\exists *AxiomVariable* • *AxiomPredicate*

- Initial state

\exists *State'* • *StateInit*

Verification with Z/Eves (III)

- Precondition calculation
 - Determines the state under which operation is successful
 - Feasibility of operation
 - Verify API robustness

$\text{pre } Operation = \exists State' \bullet Operation \setminus outputs$

$\forall State \bullet \text{pre } Operation$

- Verified refactoring
- Properties of interest

Verification with Z/Eves (IV)

- No proof obligation (PO) generation
 - Can be extended – Z/Eves uses *Z/LaTeX*
- Choose the amount of push-button proof
 - Need *housekeeping* rules to support high levels of automation
 - Extracted generic rules & theorems
 - Some theorems not-automatable – need manual declaration of proof commands

Mechanisation outcomes

- Finding missing state and precondition invariants
- Model extension with security, functional and other properties
- Reusable verified data structures
- Generic toolkit lemmas
- Refinement & code generation
- Case studies, benchmarks, verification exercises

Summary

- Tool support for mechanisation of existing Z specifications
- Z/Eves – a user friendly theorem prover for Z specifications
- Increasing repository of case studies, proofs and reusable lemmas

- Questions?